aLIGO Guardian Overview

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Introduction and Overview

Introduction

Guardian is the new aLIGO automation system.

Its job is to control the global state of the interferometer, i.e. manage and automate the locking sequence by coordinating the states of all interferometer subsystems.

It is also designed to be a useful tool for interferometer commissioning. It should aid in the commissioning process, and allow for easy reconfiguration of the interferometer to enable various commissioning tasks.

Overview

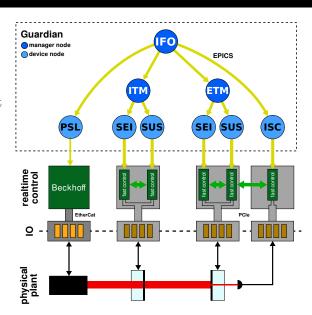
Guardian is a hierarchical state machine.

- Distributed guardian processes (**nodes**) oversee particular domains of the interferometer.
- Each node understands **states** for its domain. State code describes actions on the domain, such as changes to the plant and/or verification of the configuration.
- A hierarchy of nodes control the full IFO, with upper level **manager** nodes controlling sets of **subordinate** nodes, down to lowest level **device** nodes that talk directly to the real-time front-ends and Beckhoff.

Overview

Manager nodes control subordinates, with lowest level device nodes that talk directly to the real-time system.

A single IFO manager node will sit at the top coordinating the state of the entire interferometer.



System and state behavior

Core state machine engine daemon

The core of Guardian is the **node daemon**. It is essentially a state machine execution engine.

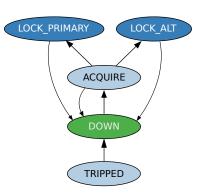
Node daemons execute **system modules** that describe the state graph that the daemon follow and the code to be executed during each state.

They run continuously, responding to system changes and user input.

Each guardian system is represented as a **directed graph**.

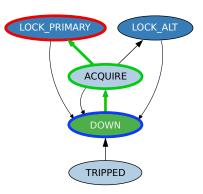
The system **states** are represented by nodes in the graph, each consisting of executable code.

Graph edges represent possible transition between states. Edges have zero content, representing only which states are accessible from each other.



Guardian accepts commands in the form of a **request state**.

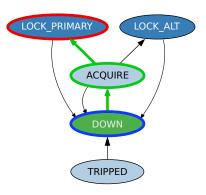
Guardian then looks at the **current state** and calculates the shortest **path** to reach the request.



Code for the current state is executed in a loop until a completion condition is returned.

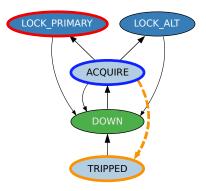
If there are more states in the path, the system transitions to the next state and immediately starts executing the new code.

If no exit condition is achieved, or the current state is the requested state, the code continues execution.



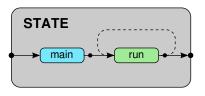
States can also specify jump transitions that bypass the normal dynamics of the graph.

Jumps represent transitions that are in some sense "undesirable" (e.g. lock loss, watchdog trip, etc.). They are ignored in normal path calculations



State behavior

The states themselves have very simple behavior. There are two state *methods* (i.e. functions).



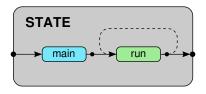
They are essentially identical except that:

main executed once immediately upon entering state.

run executed in a loop. Used to continuously check for state completion conditions.

State behavior

The states themselves have very simple behavior. There are two state *methods* (i.e. functions).

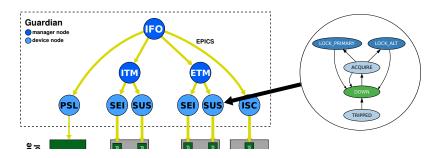


If either state returns True the state completes and guardian transitions to the next state ("edge transition").

Either method can also return the name of a state in which case guardian will **jump** immediately to that state ("jump transition").

Nodes

Each system is a **node** in the full IFO automation environment:



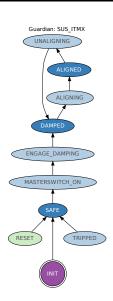
Manager nodes control subordinates by requesting states and monitoring their progress towards achieving that request.

Commissioned systems: SUS

SUS was the first subsystem automated (also the simplest). All suspensions nodes have identical guardian structure, replying upon a base SUS.py module and a sustools.py library interface.

Fully handles recovery to aligned state from watchdog trips.

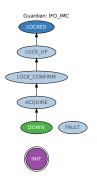
Currently (as of this week at LLO) not touching the alignment offsets. Will need to revisit how alignment offsets are managed.

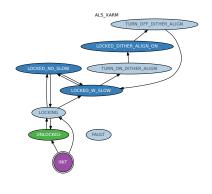


Commissioned systems: IMC and ALS

The **IMC** was next, and was the first "non-trivial" control. It is now deployed at both sites (although still being tweaked).

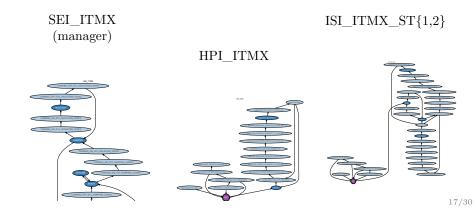
The **ALS** system under active development at LHO.



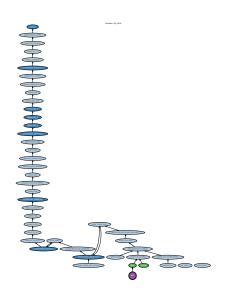


Commissioned systems: SEI

Next was **SEI** at LHO, the most complicated subsystem from an automation perspective. Fully handles recovery to full isolation, coordinating HPI and ISIs. Chamber seismic system consists of multiple nodes managed by a single "chamber manager" (HAM: 3 nodes, BSC: 4 nodes):



Under development: Full IFO Locking



Primary IFO locking node: IFO_LOCK

Full IFO lock acquisition node at LLO.

This node is also under heavy development, but starts to give a sense of what the full lock acquisition procedure will look like.

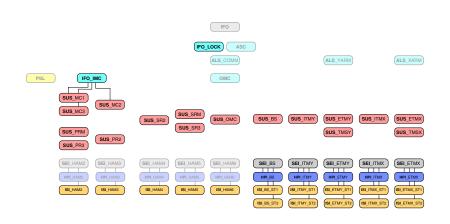
User code status

- SUS deployed at both sites.
- SEI deployed at both sites.
- IMC deployed at both sites.
- ISC is now primary focus, heavy development at LLO this summer.
- Working on improved system monitoring, front end channel monitoring infrastructure now being deployed.

Full system overview



Full node projection



Tracking Guardian status

GRD EPICS records

Guardian nodes broadcast their state and status via a set of EPICS channels, all prefixed with:

MODE	WORKER
REQUEST_S	NOTIFICATION
REQUEST_N	USERMSG
STATE_S STATE_N TARGET_S	GRDMSG ERROR
TARGET_N	LOGLEVEL
STATUS	VERSION

Primary status channel: MODE

<IFO>:GRD-<NODE>_MODE

enum record representing the execution mode of the node:

0:INIT	system initializing
1:RELOAD	reloading user code
2:STOP	worker terminated and no new code exe-
	cuted; state preserved
3:PAUSE	no new code executing but current method
	left to complete; state preserved
4:EXEC	primary code execution
5:MANAGED	primary code execution, managed mode,
	i.e. no auto-recovery from jump transi-
	tions

Primary status channel: ERROR

<IFO>:GRD-<NODE>_ERROR

enum record representing error condition in node:

0:False	no error
1:True	error present; all code execution stopped
2:Connect	channel access connection error; current
	state method is re-executed until the chan-
	nel access error clears

Primary status channels: REQUEST, STATE, TARGET

```
<IFO>: GRD-<NODE>_REQUEST_S
<IFO>: GRD-<NODE>_REQUEST_N
<IFO>: GRD-<NODE>_STATE_S
<IFO>: GRD-<NODE>_STATE_N
<IFO>: GRD-<NODE>_TARGET_S
<IFO>: GRD-<NODE>_TARGET_N
```

These channels, for the requested state, current state, and target state, are the string (S) and numeric (N) representations of **usercode states**.

States enumeration

The data acquisition "frame writer" is currently only capable of recording numbers in the LIGO frames. This means we need to enumerate all the guardian states (map to number).

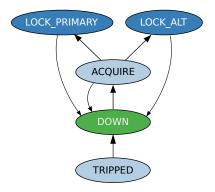
The guardutil program (part of the guardian package) is currently the best way to get state information:

guardutil

guardutil has useful functions for developing, analyzing, etc. guardian systems.

controls 0\$ guardutil graph SIMPLE_LOCK

graph draws system graphs
states list system state
enumeration
code list all usercode
print print general system info
(combo of above)
etc.



guardutil

Example run on my laptop:

```
servo:~ 0$ IFO=L1 USERAPPS=~/ligo/userapps
           quardutil states IFO_LOCK
 INIT *
 IDLE *
 DOWN *
3 TURN_ON_PR2_WFS *
4 TURN ON CHARD WFS *
 DARM_ON_ASQ *
6 CARM ON TR *
7 TURN ON AS WFS *
8 CARM ZERO OFFSET *
9 ALS COMPLETE *
10 TURN_ON_INPUT_WFS *
11 DRMI_LOCKED *
12 LOCKED *
```

States enumeration issues

Currently, guardian assigns an arbitrary number to each state at load time. As long as the code doesn't change the state enumeration doesn't change. However, if states are added or removed the enumerations will change. This means that the:

The state enumeration for a particular point in time is a function of code loaded at that time.

Not a good situation for tracking state over time. This needs to / will be improved. Currently plan: